

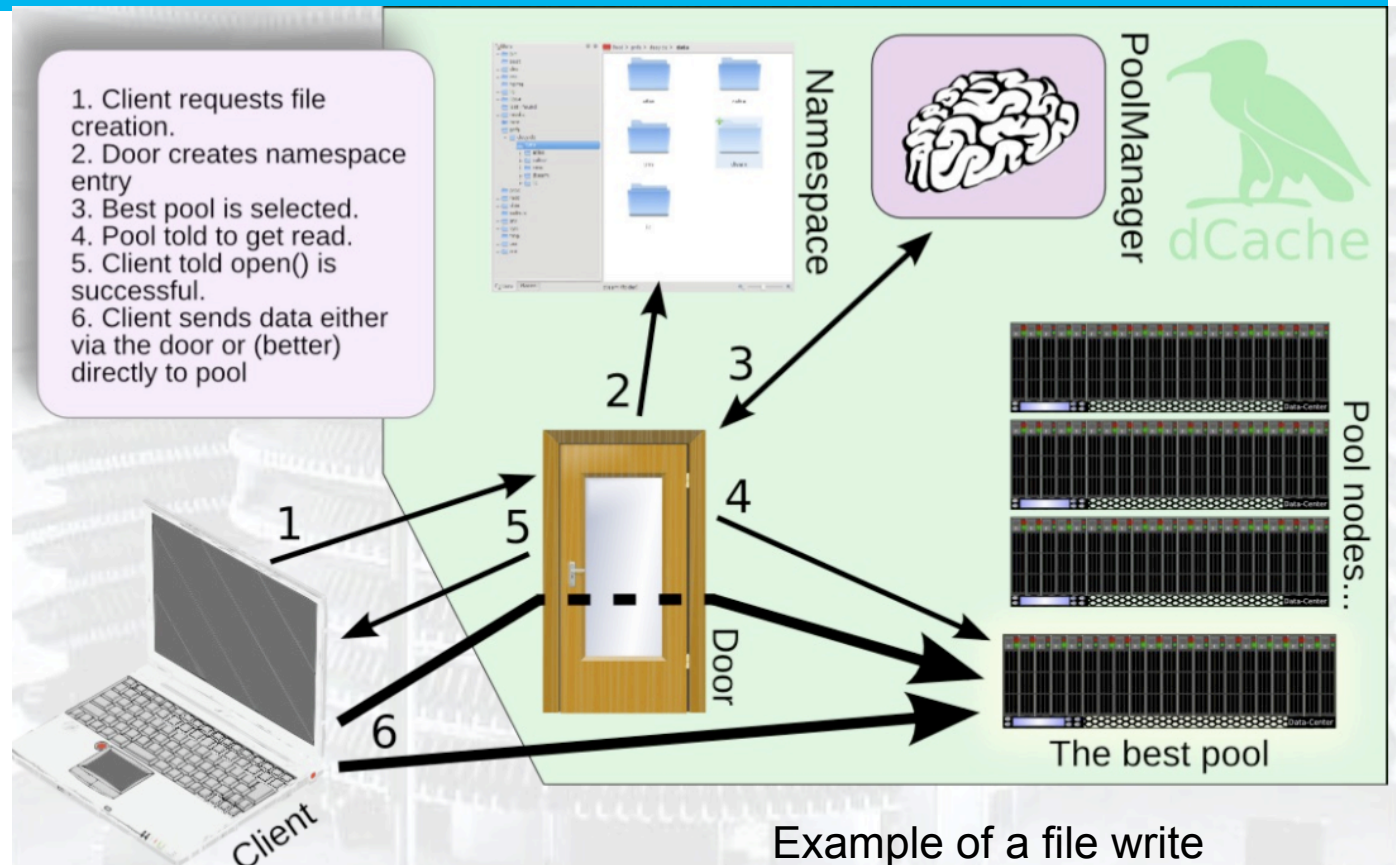
# LHC Data Analysis Using NFSv4.1 (pNFS): A Detailed Evaluation

- Short introduction into dCache and NFSv4.1 (pNFS)
- First simple tests
- ATLAS HammerCloud
- CMS Analysis
- Reading ROOT files

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Stadie (Uni Hamburg)  
Taipei, 10/20/2010 CHEP 2010

# dCache in a nutshell

- Storage system, developed at DESY, FNAL and NDGF
- Objects stored: Files
- Files in pools, pools on poolnodes, many of them
- Client connects to a door, which speaks the desired protocol
- At the end the file is transferred directly between pool and client

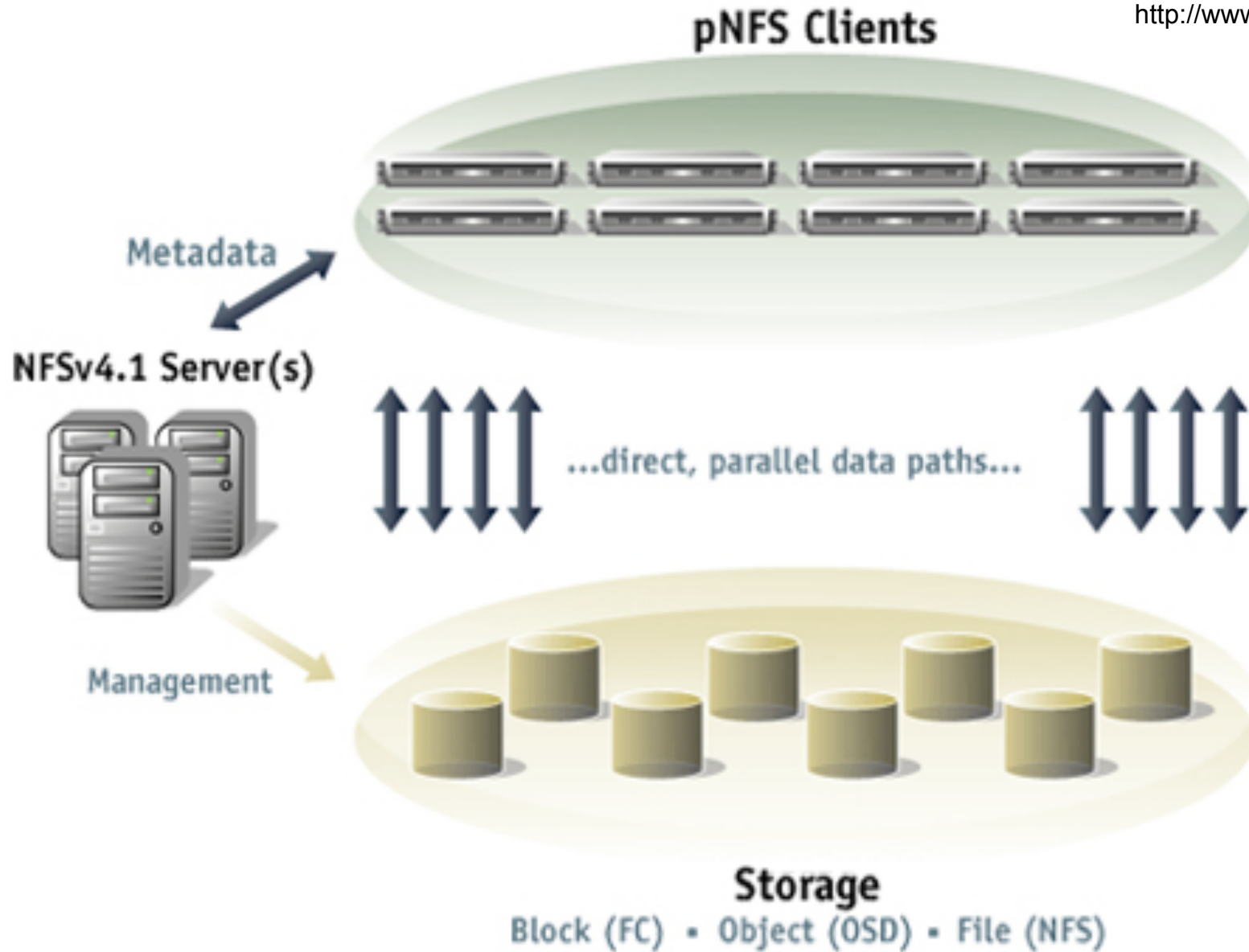


In reality a little bit more complicated  
Many talks and posters around dCache at CHEP

Check <http://www.dcache.org/>

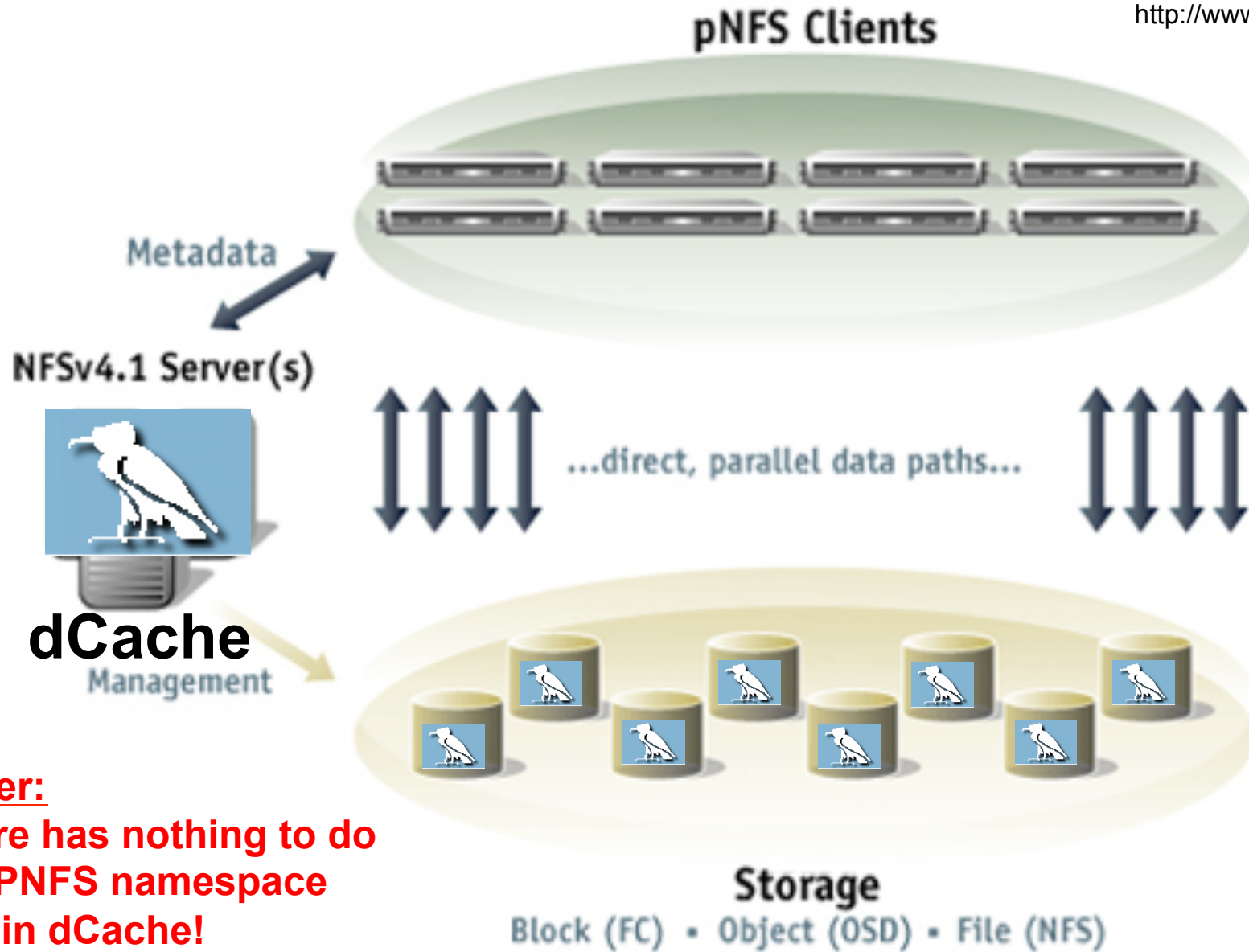
# NFS v4.1 / pNFS from the infrastructure view

<http://www.pnfs.com/>



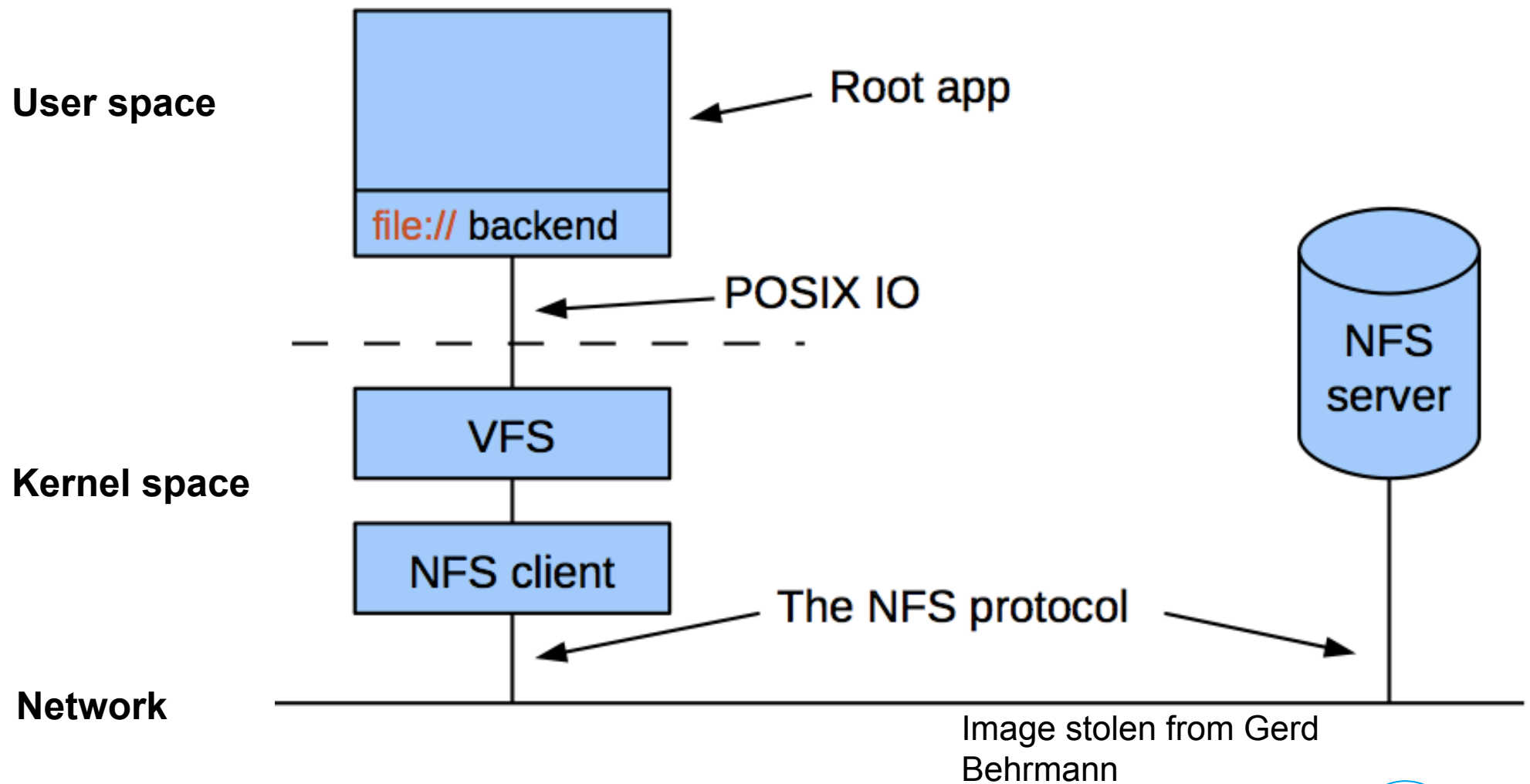
# NFS v4.1 / pNFS from the infrastructure view: adding dCache

<http://www.pnfs.com/>



**Disclaimer:**  
pNFS here has nothing to do  
with the PNFS namespace  
provider in dCache!

## ... a look from the client side



# 11 reasons why one should care about NFS 4.1

- 1) High latency link performance
  - Batching of several components, reducing number of network ops, bidirectional RPC
- 2) Proper authentication and authorization
  - Kerberos, X509 under investigation, ACL
- 3) Introduction of sessions with NFS 4.1
  - Decoupling transport from client
- 4) Parallel NFS (remember the plots to pages before)
- 5) Standardization: RFC 5661, IETF Proposed Standard
- 6) Industry backed: NetApp, Microsoft, Panasas, EMC, IBM, ...
- 7) Client availability:
  - Linux (more details later), Solaris available, Windows (U.Michigan)



# 11 reasons why one should care about NFS 4.1 (contd)

## 8) Server available:

- NetApp, IBM, Oracle, EMC, IBM,...
- dCache, DPM in WLCG context

## 9) Clients provided by industry:

- Real POSIX IO, caching provided by OS & tuned by experts, no apps modifications

## 10) Funding secured

- EMI funds NFS 4.1/pNFS in DPM and dCache, HGF (D) additional funds for dCache

## 11) Simple migration path

- Server: No data migration needed, NFSv4.1 (pNFS) is additional protocol
- Clients: user file:// -> Unifies access for dCache, DPM, GPFS+Storm

**OK, and how does the reality look like for HEP applications?**

("11 reasons" stolen from Gerd Behrmann)



# Evaluation: The testbed in the DESY GridLab

## CPU Cluster

### Batch&CE:

CREAM-CE  
glite-CREAM-3.2.6-0  
SL5.3

### Clients:

32x DELL M600 blades  
(16x in the beginning)  
2x4 cores @ 2.5 GHz  
16 GB RAM  
1 Gbit Network  
gLite-WN 3.2.7-0  
SL 5.3  
2.6.36-rc3.pnfs

### Mount on client:

`dcache-head:/pnfs on /pnfs type nfs4 (rw,minorversion=1,rsiz=32768,wsiz=32768)`

## Network

Force 10  
Gbit  
Switch

4x10 Gbit  
links to  
Arista

Arista  
10 Gbit  
Switch

## dCache Storage

1.9.10pre

### dCache Head-Node

4 core, 8 Gbyte RAM  
1 Gbit Network  
SL 5.3  
2.6.18-194.3.1.el5

### Poolnodes:

5x DELL R510  
2x4 cores @ 2.27 GHz  
12 GB RAM  
10 Gbit Network (Intel)  
SL 5.3  
2.6.18-194.3.1.el5  
2x2 TB RAID-1 System  
2x10 TB RAID-6 Data



# What to expect from testbed?

- > Maximum BW from **server → Clients**: 40 Gbit (link between two switches)
- > Maximum BW from **one pool → Clients** (alone): Theoretical 10 Gbit
  - Measured to 5.6 Gbit/s using iperf
- > Maximum BW from **Disk RAID → local /dev/null**
  - Measured between 520 MByte/s (few streams) and ~300 MBytes/s (random read)
- > **So, maximum bandwidth from Server-Disks → Network → Client /dev/null**
  - Something between 1.5 GByte and 2.5 GByte/s
  - 32x1-Gbit clients can saturate this
- > CPU ~ 1/2 Tier-2 whereas Storage ~1/4 Tier-2
  - Clients able to really stress the storage system
  - **Storage undersized (on purpose!)**



# First simple test

## > Simple I/O

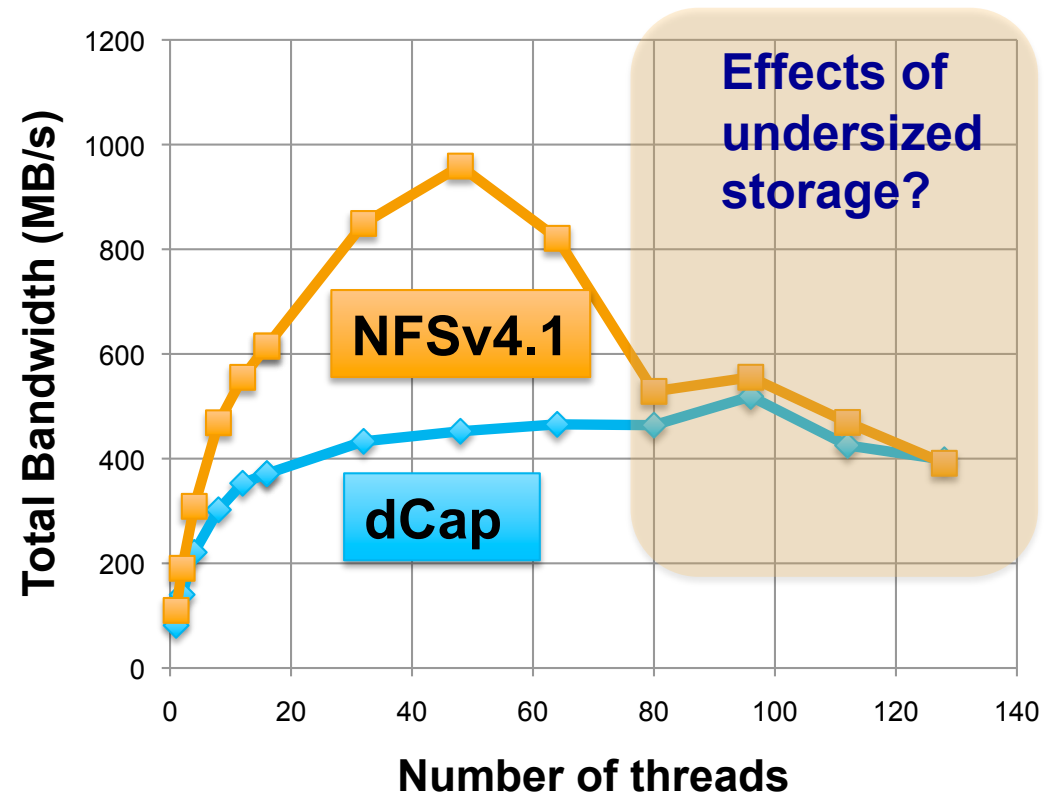
- Reading file to /dev/null
- No caching (read once, not jumping around in file)
- A maximum of 128 clients (16 nodes)

> NFS behaves better than dCap up to a certain limit

> We have no definite answer for this effect, suppose congestion on the server

- Probably due to undersized storage

→ **Needs further investigation**



# Stability tests

- > Untaring the Linux Kernel into NFS 4.1 ✓
  - up to 16 parallel jobs (only 16 clients)
  - Works, slowly, but no problems observed with recent kernels
- > CFEL Production Transfers from SLAC to DESY ✓
  - 13 TBytes over 10 days
  - 100 GBytes average file size
  - No crash
- > High-Latency test: “recursive ls -l” 60k files over DSL from home ✓
  - Slow, but works
- > 128 clients simultaneously writing into same file (by mistake) ✗
  - Client nodes got stuck
  - Server OK
- > Clients got stuck once during ROOT tests, needed reboot ✗



# ATLAS HammerCloud test: The setup

## > The Data:

- Official ATLAS MC samples (7 TeV, preferably no minbias, few jets)
- AODs, reconstructed with athena 15.6.8
- 33 TB data in total

## > The Analysis

- standard AOD analysis reading Trigger and many Muon variable
- Athena 15.6.6, ROOT 5.22/00h (no treecache reading used)

## > Initial difficulties:

- CREAM-CE not visible, neither in Information System, nor “in the Cloud”
- dCache not a fully Grid-SE, had to provide file lists as input

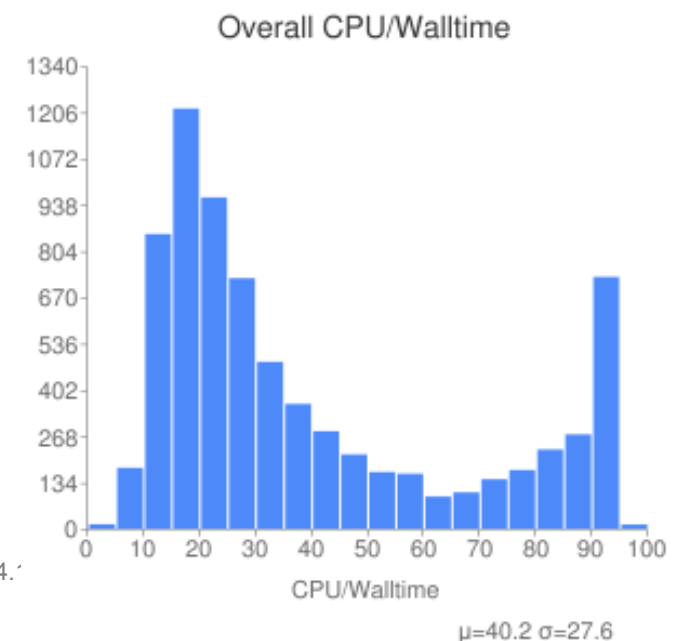
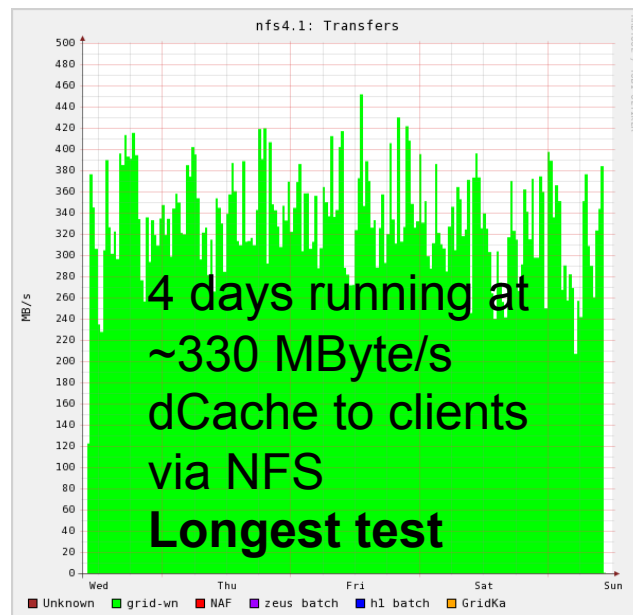
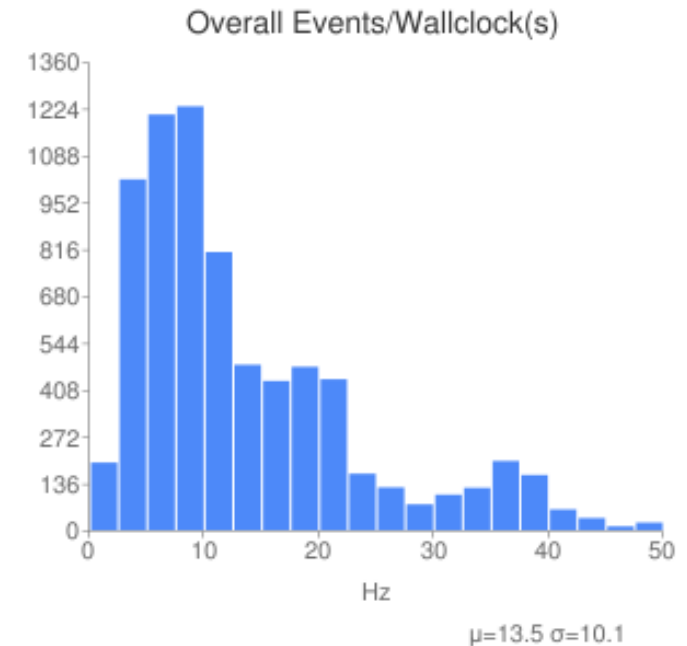
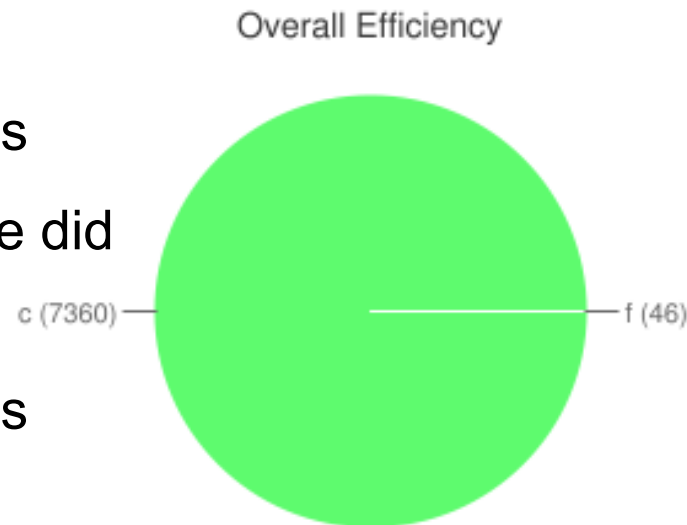
## > More on HammerCloud

- This is the standard ATLAS application to test the performance of sites
- Parallel session 36, Dan van der Ster
- Poster PO-MON-036, Federica Legger



# ATLAS HammerCloud test: The results

- > 8248 jobs in total
- > Cancelled after 4 days
- > Longest single test we did
  - No trouble during test
- > Reasonable outcomes (events/s,...)
- > No comparison made to dCap (yet)



# CMS Analysis: Setup

- > Job submission done via the Grid and *grid-control*
  - Ability to freely define CE (which was “hidden” in our case)
  - Make use of “private” SE: Custom manipulation of the CMS Trivial File Catalogue
  - <https://ekptrac.physik.uni-karlsruhe.de/trac/grid-control>
- > Muon analysis. Dataset: 1.7 TB in 308 RECO files
- > Executable: filetest is stripping into PAT Ntuple out of the CMSSW framework
  - Using 5.22 ROOT version shipped with CMSSW
- > One typical use-case on the DESY National Analysis Facility
- > Not much CPU, nearly only I/O
- > Evaluation of performance metrics in CMSSW framework job report (Andrzej Wronka (summerstudent at DESY))



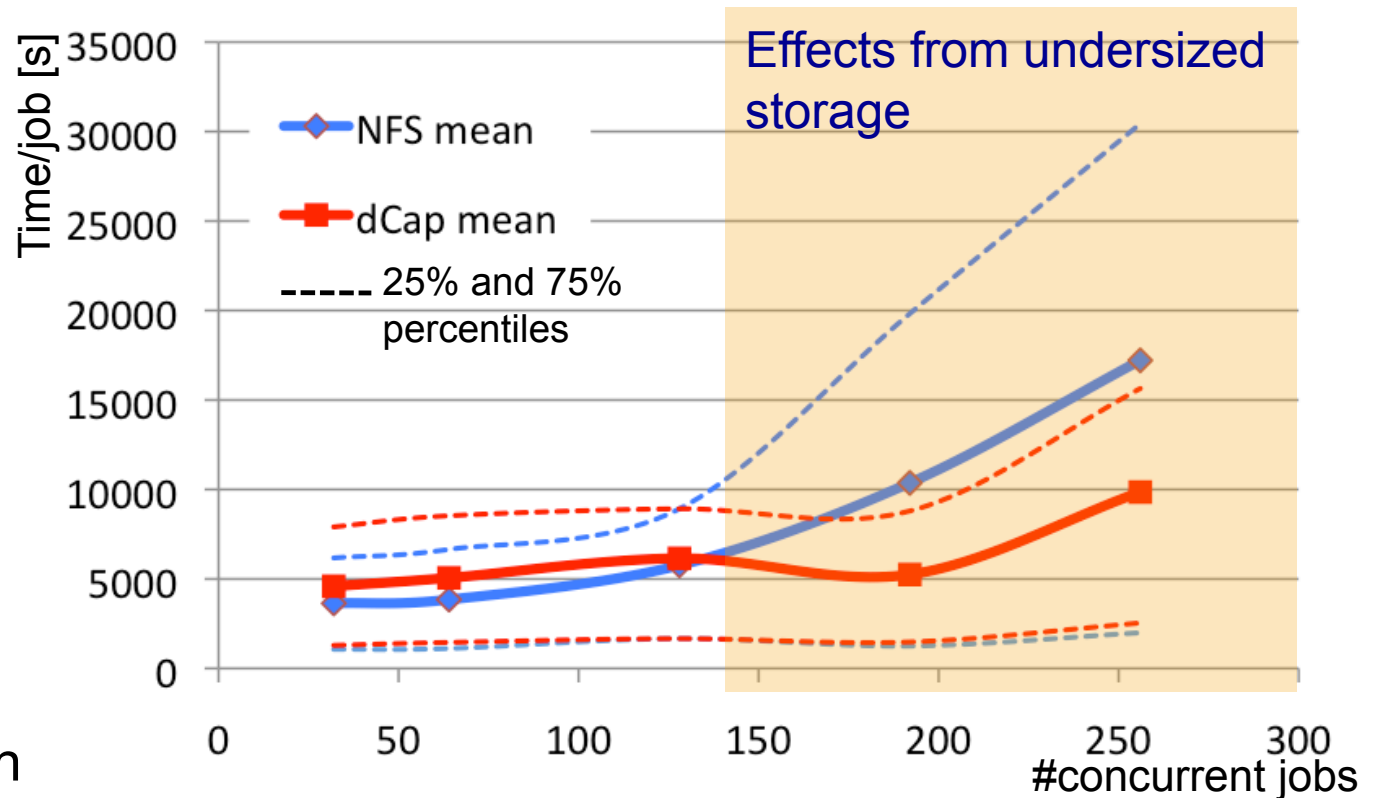
# CMS Analysis: Results

## Below ~128 jobs:

- > NFS 20% faster than dcap

## Above ~128 jobs:

- > NFS performance degrades, dcap only slightly degrades
- > Not yet fully understood, suspect numbers of threads in dCache NFS server
- > Checked that client congestion not fault

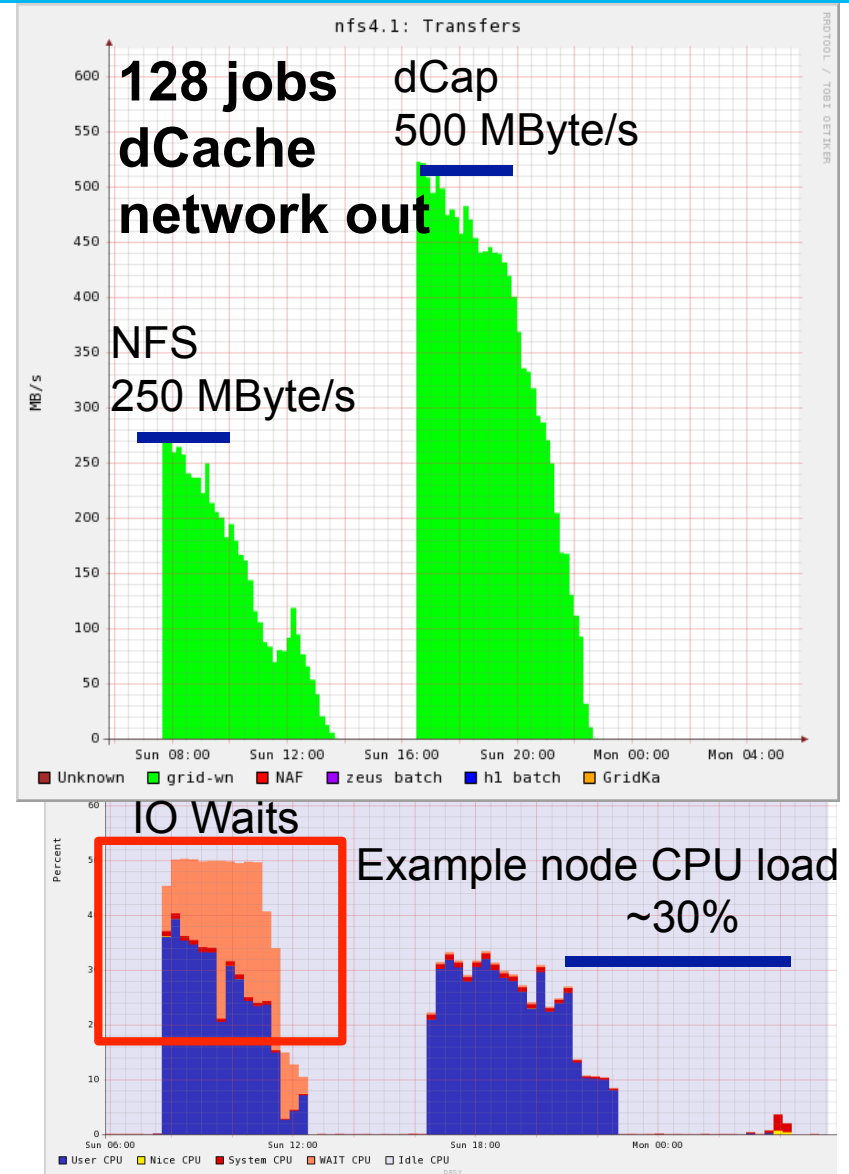
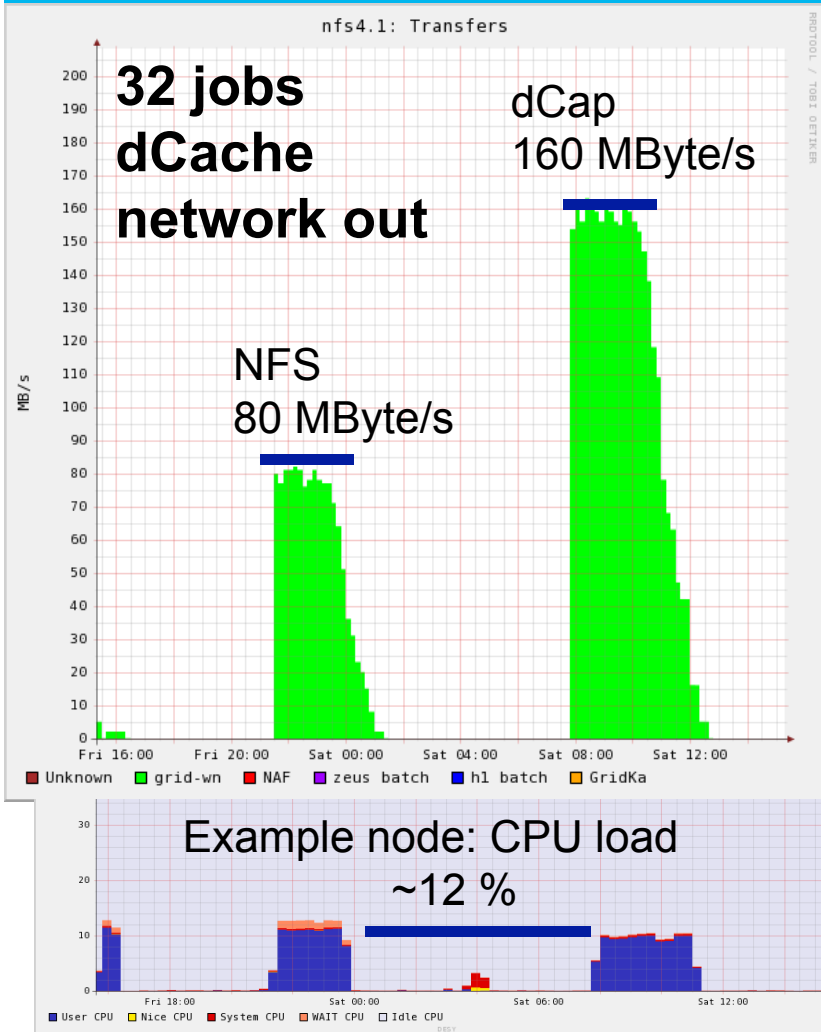


## Effects of File system cache:

- > dCap reads 2.5 times more data than NFSv4.1 (dCache billing logs and network monitoring plots): Next slide:



# CMS Tests: A look at dCache and one node



- > IO waits gets more important for NFS at higher numbers of concurrent jobs
- > Less network traffic for NFS

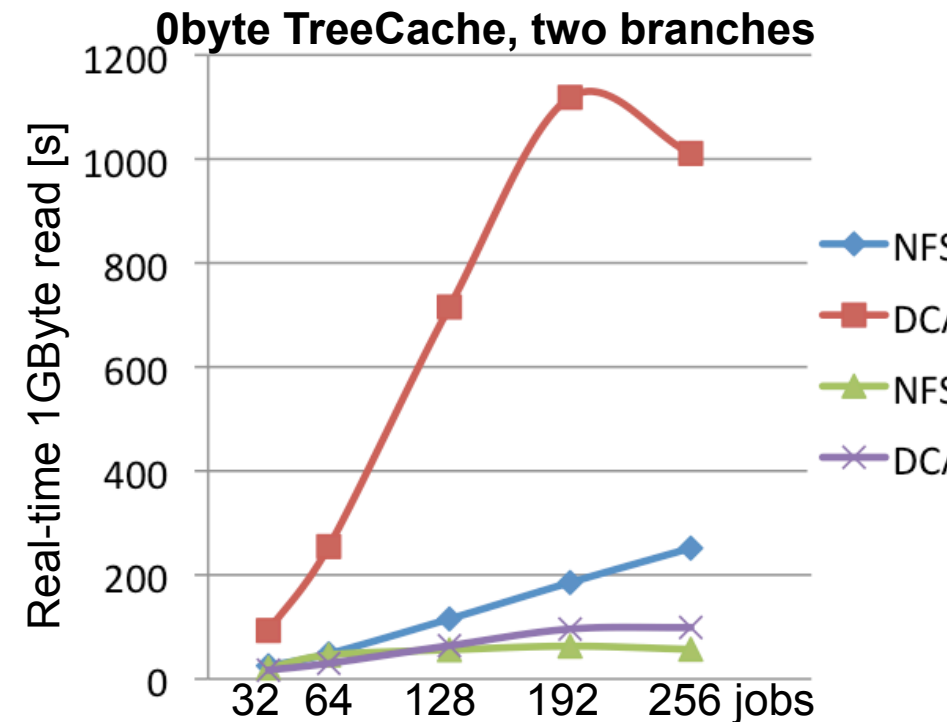
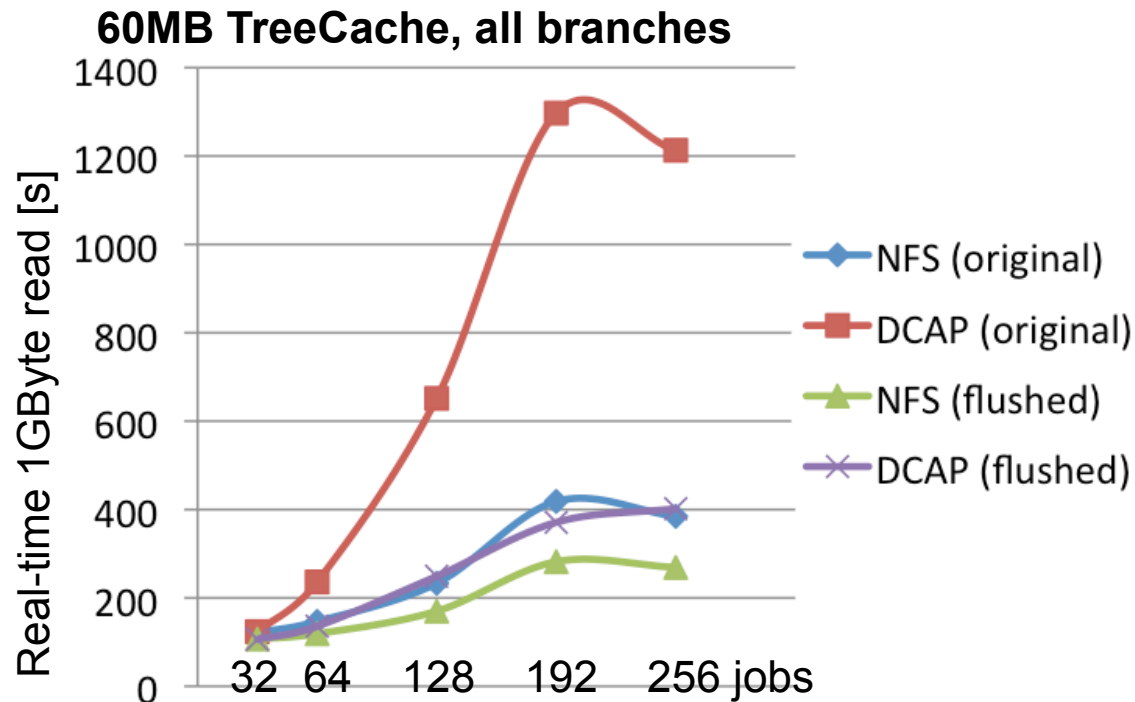


# Half-Synthetic ROOT tests: Setup

- > New ROOT version 5.27.06, compiled with dCap support
- > Files provided by René Brun: `atlasFlushed.root` (re-organized files with optimized buffers) and `AOD.067184.big.pool_4.root` (some other original file) (flushed: 1GByte, original 1.3 GByte)
- > Test script provided by René: simple script reading events: `taodr.C`
- > Different test runs:
  - Reading via NFS or dCap
  - Reading with 60MByte TreeCache, or with 0Byte TreeCache
  - Reading all branches or only 2 branches
  - 32, 64, 128, 192 or 256 jobs running in parallel
- > Last minute-result! Have not spoken with ROOT people!



# Half-Synthetic ROOT tests: Results



- NFS better for original and flushed files than dCap
  - Flushed: not much difference, original: Large difference
- TreeCache helps, NFS adds additional speed
- Peak at 192 clients not understood
- Remember: Just going through events and doing nothing ... not really representative for analysis

# Kernel availability

- Kernel used for evaluation : 2.6.36\_rc3
- NFS 4.1 (pNFS) kernels expected in SL6.(>2)
- 2.6.36 back-port to SL5 available from DESY
  - Plus 'mount tools' RPM.
  - Kernel will very likely not cover all hardware setups.
- With a Joined Effort (e.g. CERN, FNAL, DESY), we would be able to provide an SL5 with NFS 4.1 (pNFS) kernel within months. (If we really want)

**Patrick Fuhrmann**  
**@ GDB 10/13/2010**



# Summary

## > Set up different use cases

- Synthetic, ATLAS HammerCloud, CMS analysis, ROOT files
- No change to experiments applications needed
- Managed to be run and steered by non-experts (like me)

## > Set up a test bed comparable to a small Grid site

- **Underpowered w.r.t dCache storage:** Able to see bottlenecks

## > Presented results

- **Synthetic:** Provide general performance and stability measurements of NFS 4.1/pNFS
- **ATLAS HammerCloud:** Stable and well-performing running over four days
- **CMS analysis:** See effects of FS cache, excellent behavior of NFS up to some point
- **ROOT files:** See effects of FS cache, better performance than dcap, even with most recent ROOT version and with TreeCache enabled

## > NFS 4.1/pNFS has advantages over traditional proprietary protocols

## > We now know: Performance is one of them!



# Future

- > More tests needs to be done, some issues have to be understood and fixed
- > Remember: NFS4.1 (pNFS) is not dCache only. NetApp have promised to give us a test storage a.s.a.p. (unfortunately not in CHEP timeline...)
  - DPM: Talk by Ricardo Rocha in Parallel Session 15
- > No mentioning of security, authentication, authorization here. This needs to come next (and will!)
- > Maybe it is time to think about a backport of NFS 4.1 (pNFS) into SL5 kernel? Could this be a combined effort? Would be a temporary effort!



# Backup 1: Complete set of ROOT result plots

